

1 **III. Organization of Binary Logic (Set Theory) and Trial of csQCA**

2 **III-1. Content of This Chapter**

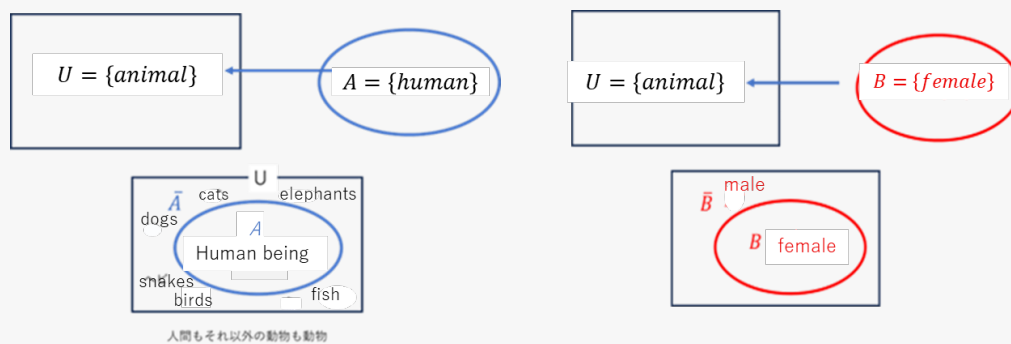
3 Using the dataset targeted for numerical analysis, we will verify Lipset's hypothesis in
4 a set-theoretic manner. This verification method has already been described in the
5 book by B. Rihoux and C. Ragin (2009) and is not original to the commentator.

6 However, to deepen understanding, basic explanations have been added. Here, we will
7 explain csQCA: crisp set qualitative comparative analysis, and its purpose is to
8 organize the binary logic that forms the basis of QCA.

9

10 **III-2. Set Theory, Boolean Operations, and Truth Tables**

11 When we say, “Humans are animals,” from a set-theoretic perspective, it means that
 12 “the set of humans is within the set of animals.” Conventionally, the set representing
 13 the whole is often denoted as U (probably derived from Universe). The set of humans is
 14 denoted as A . In logic, a statement like “Humans are animals” is called a proposition.
 15 This term is used because it can be questioned whether humans are animals (true) or
 16 not (false). This proposition can be expressed in set notation as $A \subseteq U$. Since there are
 17 many animals other than humans, it might be more accurate to write $A \subset U$. This
 18 categorizes animals by species, but animals can also be categorized by reproductive
 19 ability. Thus, a proposition like “Animals that can produce offspring are called females”
 20 can be formed. If we denote the set of females as B , then $B \subset U$. A diagram like Figure
 21 6, which expresses such set-theoretic relationships, is called a Venn diagram.

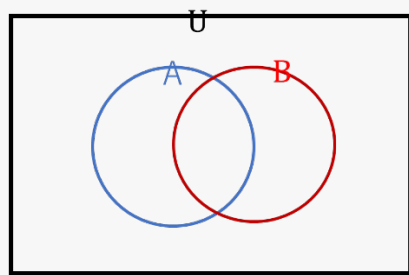


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23 Fig.6 Venn diagrams of propositions of “human beings are animal” and Animals that
 24 can produce offspring are called females”

25 Propositions create subsets A and B within a set U (in this case, animals) and
 26 simultaneously create sets of things that are not A and not B . In other words, they
 27 create sets that are false with respect to the proposition. These are called the
 28 complements of A and B . Within the whole set U , the part that is not A is called the
 29 complement of A . In set notation, complements are denoted as \bar{A} or A^c and \bar{B} or B^c .
 30 When simplifying further, they are represented by lowercase letters a and b . The
 31 relationship between A and \bar{A} is called a complementary relationship.

32 Figure 7 illustrates the world of the animal kingdom (U) by combining Proposition A:
 33 “Humans are animals.” and Proposition B: “The animals that produce offspring are
 34 female.”



$$A \cup B = \overline{\overline{A} \cap \overline{B}}$$

De Morgan's laws



$$\overline{A \cap B} = \overline{A} \cup \overline{B}$$

Male animal except men
 \ni the neighbor dog,
 Snoopy etc.



$$A \cap \overline{B} = \overline{\overline{A} \cup B}$$

Men \ni Prof. Yagi, Sakai,
 Li, Ryan, Ted. Shiga, etc.



$$\overline{A} \cap B = \overline{A \cup \overline{B}}$$

Women \ni Mina, Georgina,
 Uwai-san etc.



$$\overline{A} \cap \overline{B} = \overline{A \cup B}$$

Female animal except huma being
 \ni doe (female deer), Kitty -
 chan, etc.

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Fig 7. Structure of animal world composed from two preposition

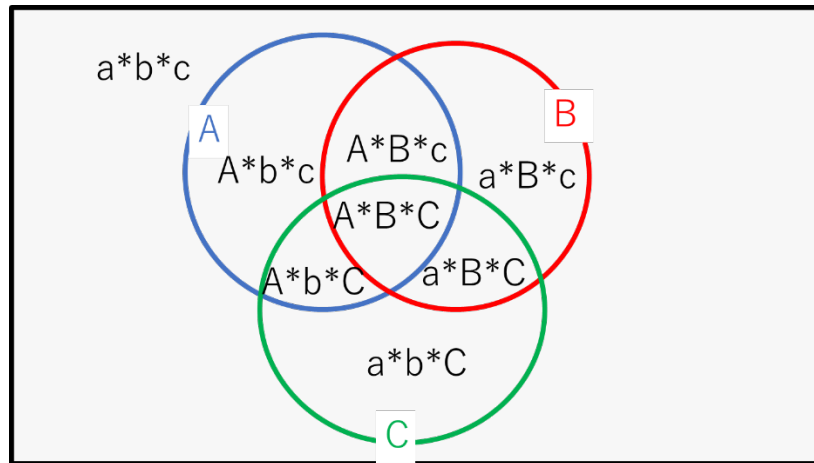


Fig 8. Venn diagram when condition A,B, and C are combined

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When two sets are combined, a set that belongs to both sets simultaneously is formed. This is called a product set. Additionally, a set that belongs to either or both of the two sets can also be formed. This is called a sum set. In set operation notation, the product set is represented as $A \cap B$ and the union set is represented as $A \cup B$. As shown on the right side of Figure 6, the entire set U (in this example, all animals) is completely filled with four subsets. In set operation notation, it can be written as follows. That is, as subsets of U, there are four intersection sets:

47

$$\overline{A} \cap \overline{B} \cup A \cap \overline{B} \cup A \cap B \cup \overline{A} \cap B = U$$

48

The sum of these four intersection sets is U. When described with two binary characteristic values such as type and gender, U has 2^2 subsets of product sets. Figure 8 shows the subsets in the form of a Venn diagram when three conditions A, B, and C are combined. Here, the product sets are described using Boolean operation notation, and the complement sets are represented in lowercase. When three characteristics are

52

53 combined, U has 2^3 subsets of product sets. When described with k characteristics, U is
54 filled with 2^k subsets of intersection sets (different combination of states or conditions).

55 When considering characteristics that distinguish members of a set, such as being
56 human or female, as conditions or states that the animal satisfies, these conditions can
57 be expressed in binary terms: true (1) when the condition is met and false (0) when it is
58 not. When there are two conditions, the logical conjunction (AND) refers to whether
59 both conditions are met. If both are met, the logical conjunction is true (1); otherwise, it
60 is false (0). This corresponds to the product set (intersection of sets). The logical
61 disjunction (OR) refers to whether at least one of the two conditions is met. If either
62 condition is met, the logical disjunction is true (1); if neither is met, it is false (0). This
63 corresponds to the sum set (union of sets).

64 For example, let A represent “being human” as true (1) and its negation a: “not being
65 human” as false (0), and let B represent “being female” as true (1) and its negation
66 b: “not being female” as false (0). The logical conjunction, “being human being and
67 female” is true (1) when both “being human” and “being female” are true, and false (0)
68 for all other combinations. Thus, women are true, represented by the logical
69 conjunction as 1, while all other males and animals are false, represented as 0. The
70 logical disjunction “being human” or “being female” is true (1) if either “being human”
71 or “being female” is satisfied. Therefore, the set of male animals excluding men has a
72 logical disjunction of 0, while all others have a logical disjunction of 1. The logical
73 conjunction is presented as $A \wedge B$, and the logical disjunction as $A \vee B$.

74 Logic is the condition that distinguishes the set from others and is the boundary line in
75 a Venn diagram. If the condition distinguishing A from others is (α, β) and the
76 condition distinguishing B from others is (β, γ) , then

$$77 \quad A \wedge B \rightarrow (\alpha, \beta, \gamma), \quad A \vee B \rightarrow (\beta, \gamma)$$

78 Considering sets, if all elements (members) included in the set are represented by a, b,
79 c:

80 If

$$81 \quad A = \{a, b\}, \quad B = \{b, c\}$$

82 then

$$83 \quad A \cap B = \{b\}, \quad A \cup B = \{a, b, c\}$$

84 It might be confusing, but it can be understood by thinking, “The more conditions there
85 are, the fewer the number of suitable members.”

86 Boolean operations (Boolean logic) are methods of calculation for sets and logic, but
87 they yield the same results whether applied to sets or logic. In practical terms, logical
88 expressions can be difficult to read. Therefore, in this explanation, unless it is
89 necessary to emphasize that it is a logical operation, the notation for sets is used.

90 **Some Laws of Boolean Operations**

91 **Notation:**

- 92 • Union(sum): +
- 93 • Intersection(product): *
- 94 • Complement: \tilde{A}, A^c, a
- 95 • Members: $A = \{\alpha, \beta, \gamma\}$ α, β, γ are elements that constitute A.

96 $A \ni \alpha. \quad A \not\ni \delta$ α is included in the members. δ is not included in
97 members

- 98 • Empty set: \emptyset a set with no members

99 **Main Laws:**

100 **1. Commutative Law:**

- 101 $\circ \quad A+B=B+A$
- 102 $\circ \quad A*B=B*A$

103 **2. Associative Law:**

- 104 $\circ \quad (A+B)+C=A+(B+C)$
- 105 $\circ \quad (A*B)*C=A*(B*C)$

106 **3. Distributive Law:**

- 107 $\circ \quad A*(B+C)=A*B+A*C$

108 **4. Identity Law:**

- 109 $\circ \quad A+1=1$ The union of a part and the whole is the whole
- 110 $\circ \quad A*1=A$ The intersection of a part and the whole is the part
- 111 $\circ \quad A+A=A$
- 112 $\circ \quad A*A=A$

113 **5. Complement Law:**

- 114 $\circ \quad A+\tilde{A}=1$
- 115 $\circ \quad A*\tilde{A}=0$

116 **6. Absorption Law:**

Truth table				Result: product set and Union					
ID	condition			result		condition		result	
	A	B	C			A	B	product	union
A*B*C	1	1	1						
A*B*c	1	1	0		A-B	1	1	1	1
A*b*C	1	0	1		A-b	1	0	0	1
A*b*c	1	0	0		a-B	0	1	0	1
a*B*C	1	1	1		a-b	0	0	0	0
a*B*c	1	1	0						
a*b*C	1	0	1						
a*b*c	0	0	0						

Fig. 10 An example of procedure to make truth table

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Next, the commentator will explain the truth table. A truth table is a table that represents whether certain binary conditions are met or not using 0 and 1. The leftmost column in figure 10 contains the code IDs. For example, in the data of European countries during the interwar period, which is the subject of this analysis, the IDs of each country are recorded. In Table 10, the Boolean expression for the intersection set is included as the ID. To the right of that, the truth values corresponding to each ID are recorded in binary (0 or 1) under the column names of the conditions. Further to the right, there is a result column that shows what happens or what happened under those conditions. In the example of interwar Europe, the data is binary, indicating whether the country maintained democracy (1) or democracy collapsed (0). As an example, when conditions A and B are the values in the condition columns, the logical conjunction and disjunction results are shown.

154 **III-3. How to Use Truth Tables**

155 Truth tables are very convenient once you learn how to use them. By using Excel's
 156 sorting function to rearrange the truth table, you can achieve various things. The
 157 usage of the truth table is shown in Figure 11. In the truth table shown on the left side
 158 of the figure, the condition that leads to the result R being true (1) is only the
 159 combination of conditions $A \wedge B \wedge C$. It can be concluded that if $A \wedge B \wedge C$, then the result R
 160 is obtained ($A \wedge B \wedge C \rightarrow R$ R is a necessary condition for $A \wedge B \wedge C$). In the middle example,
 161 since the result R=1 is obtained with both $A \wedge B \wedge C$ and $A \wedge B \wedge \bar{c}$, taking their logical or
 162 gives $A \wedge B \rightarrow R = 1$. In Boolean algebra, this can be written as follows:

163
$$A \wedge B \wedge C + A \wedge B \wedge \bar{c} \rightarrow R$$

 164
$$A \wedge B \wedge (C \vee \bar{C}) \rightarrow R$$

 165
$$A \wedge B \rightarrow R$$

166 In the example on the far right, We should reach the conclusion in following route.

- 167 1. R=1 is obtained with logical disjunction of
 168 $(A \wedge B \wedge C) \vee (A \wedge b \vee C) \vee (a \wedge B \wedge c) \vee (a \wedge b \wedge c)$.
 169 2. There is no consistency in truth value in condition A and B and there is a
 170 consistency in truth value in condition C
 171 3. It can be judged that if C=1, necessarily R=1.

set	conditions			Result
	A	B	C	R
A*B*C	1	1	1	1
A*B*c	1	1	0	0
A*b*C	1	0	1	0
A*b*c	1	0	0	0
a*B*C	0	1	1	0
a*B*c	0	1	0	0
a*b*C	0	0	1	0
a*b*c	0	0	0	0

set	conditions			Result
	A	B	C	R
A*B*C	1	1	1	1
A*B*c	1	1	0	1
A*b*C	1	0	1	0
A*b*c	1	0	0	0
a*B*C	0	1	1	0
a*B*c	0	1	0	0
a*b*C	0	0	1	0
a*b*c	0	0	0	0

set	conditions			Result
	A	B	C	R
A*B*C	1	1	1	1
A*B*c	1	1	0	0
A*b*C	1	0	1	1
A*b*c	1	0	0	0
a*B*C	0	1	1	1
a*B*c	0	1	0	0
a*b*C	0	0	1	1
a*b*c	0	0	0	0

Boolean operation complement

$A + a = 1$

$B + b = 1$

$C + c = 1$

Contradiction

$A * B * C \rightarrow R$

$A * B * C + A * B * \bar{c}$
 $= A * B * (C + \bar{c})$
 $= A * B$
 $A * B \rightarrow R$

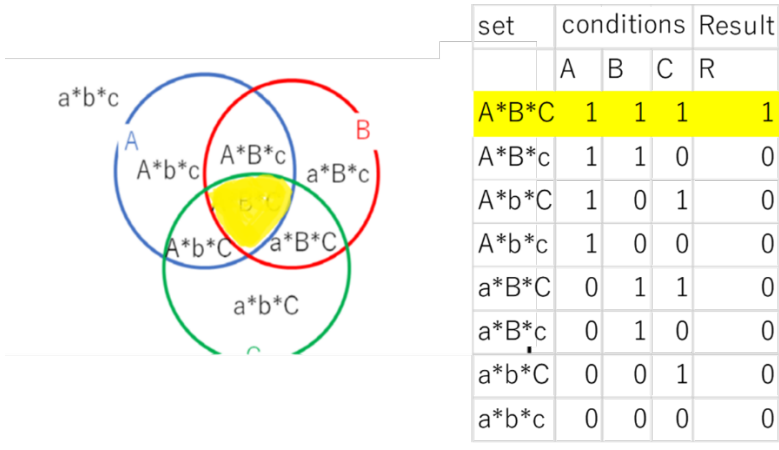
Contradiction

$A * B * C + A * b * C + a * B * C + a * b * C$
 $= A * C * (B + b) + a * C * (B + b)$
 $= A * C + a * C = C * (A + a) = C$
 $C \rightarrow R$

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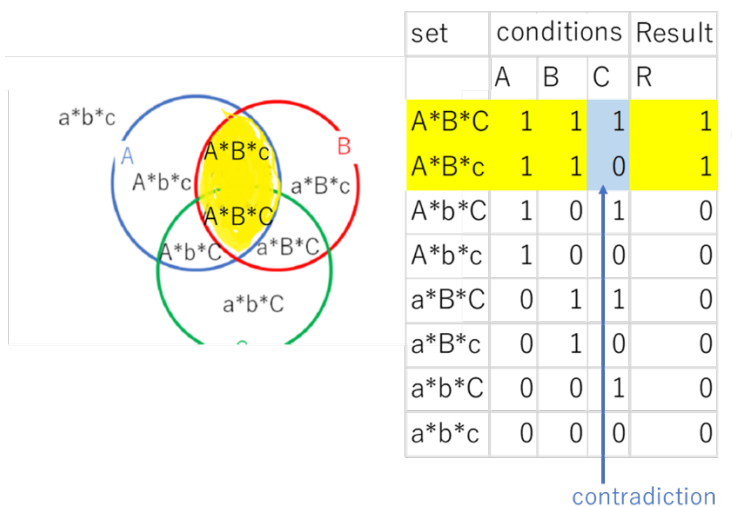
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Fig. 11. How to use truth table



$$A \wedge B \wedge C \rightarrow R$$

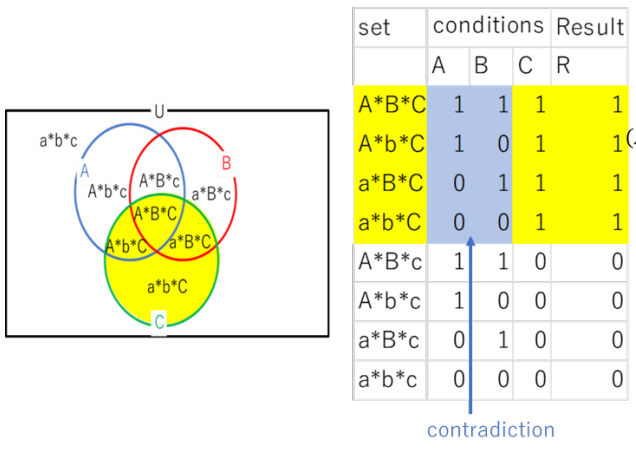
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$$(A \wedge B \wedge C) \vee (A \wedge B \wedge c) = A \wedge B \wedge (C \vee c) = A \wedge B$$

$$A \wedge B \rightarrow R$$

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$$(A \wedge B \wedge C) \vee (A \wedge b \wedge C) \vee (a \wedge B \wedge C) \vee (a \wedge b \wedge C) = A \wedge C \wedge (B \vee b) \vee a \wedge C \wedge (B \vee b) = (A \wedge C) \vee (a \wedge C) = C \wedge (A \vee a) = C$$

$$C \rightarrow R$$

Venn's diagram Truth table Boolean operation

176
177 **Fig. 12. Comparison of analyses by Venn diagram, truth table and Boolean**
178 **operation**

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180 It can be concluded, at first glance, that since only C is common, $C=1$ is enough
181 condition for R ($C \rightarrow R$) using truth table.

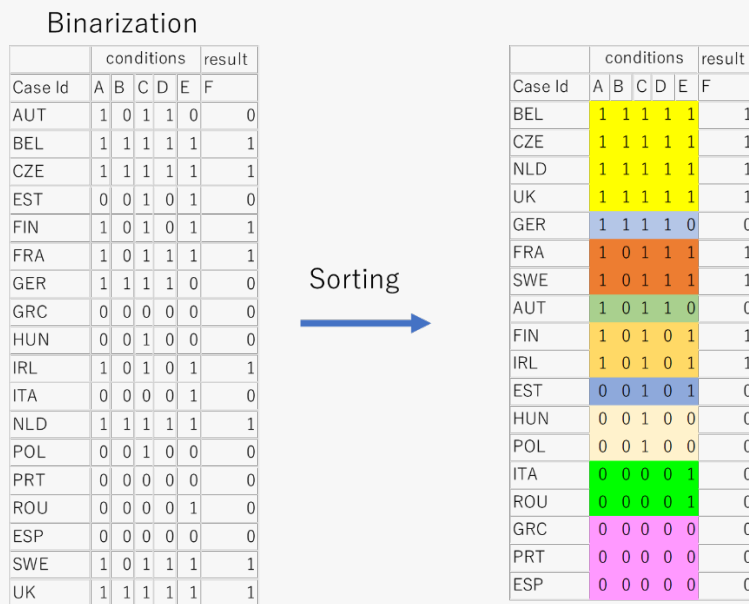
182 In Figure 12, the analysis of the truth table from Figure 11 is shown side by side with
183 the analysis using Venn diagrams and Boolean operations. Although the
184 representations look different, the content of what is being done is the same. Any
185 method that is easy to understand and minimizes errors is acceptable, so Boolean
186 operations are not particularly meaningful at this point. However, if someone says that
187 they were added just to give some authority, they wouldn't be entirely wrong, but they
188 are not completely meaningless either. This will become clear later. At this point, what
189 is more important is how to use the Excel sheet. In the example table, there are only
190 three conditions, so the truth table has only eight rows. When there are more
191 conditions, it becomes difficult to find the rows where $R = 1$. In such cases, using
192 Excel's sorting feature to prioritize the R column makes the task easier and reduces
193 errors. There is also a QCR package for R. I took a quick look, but it seemed that the
194 meanings of the individual functions were not clearly defined. With some technique
195 and knowledge, QCR, including fsQCR, can be executed in Excel, so to understand it
196 properly, it is advisable to try executing QCR in Excel once.

197

198 III-4. Trial of csQCA Using Lipset’s Theory Verification as a Subject

199 csQCA is an analysis of binary data expressed in 0-1. Combination of explanatory
 200 variables is used to explain the dependent variable similarly to multiple regression
 201 analysis. Both the explanatory variables and the dependent variable are binary, and
 202 the results are also binary, indicating whether they can be explained or not. The data
 203 is strictly (or crisply) written in binary, and the results are also binary, hence the name
 204 “crisp set QCA.” Not only the dataset but also the conclusions are crisp. The specific
 205 procedure for csQCA has already been explained in the use of the truth table. Here, we
 206 will verify Lipset’s hypothesis, which was the subject of numerical analysis in the
 207 previous chapter. In csQCA, the data (Table1) must be binary. Therefore, thresholds
 208 are set for each data item to binarize them. Figure 13 shows the first step of csQCA:
 209 binarization and sorting. The data is from interwar Europe, which was the subject of
 210 analysis in Chapter 1. The set thresholds are A: wealth 600, B: urbanization 50.0, C:
 211 education 75.0, D: industrialization 30.0, E: political stability 9.9 (government
 212 turnover), R: maintenance of democracy 0.

213 The table on the left of figure 13 is the binarized dataset, and by using Excel’s sorting
 214 function to prioritize the columns from E to A in descending order, the table on the
 215 right is obtained. Countries with the same conditions in the truth table are grouped
 216 together and color-coded. From the top, the group with 11111 includes [Belgium](#),
 217 [Czechoslovakia](#), [the Netherlands](#), and [the United Kingdom](#), which are countries that



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Fig. 13. First step of csQCA: binarization and sorting

220 maintained democracy. The second group, 11110, includes **Germany**, a country where
 221 democracy collapsed. Following are 10111: **France, Sweden**; 10110: **Austria**; 10101:
 222 **Finland, Ireland**; 00110: **Estonia**; 00100: **Hungary, Poland**; 00001: **Italy, Romania**;
 223 00000: **Greece, Portugal, Spain**. It is noteworthy that no group includes both countries
 224 that maintained democracy and those where democracy collapsed. Theoretically, there
 225 are $2^5 = 32$ combinations of conditions. Table 10 includes the names of each country
 226 for all combinations in the truth table. The consistency in Table 10 is the proportion of
 227 countries with same outcomes. If there is low consistency in the results, that
 228 combination of the conditions cannot be used as condition to explain the results. The \emptyset
 229 in the table represents an empty set, meaning there were no countries belonging to
 230 that condition. In logical terms, an empty set is called a logical remainder. Of the
 231 theoretically possible combinations, only nine cases were observed. Of course, since
 232 there are only data for 18 countries, this is unavoidable. However, it seems that
 233 Denmark, Switzerland, Norway, etc., should be included in the analysis. It is
 234 interesting to see whether consistency is maintained even with such a dataset.

235

236 Table 10. Truth table of all combination of conditions with countries belonging the
 237 condition

Truth table

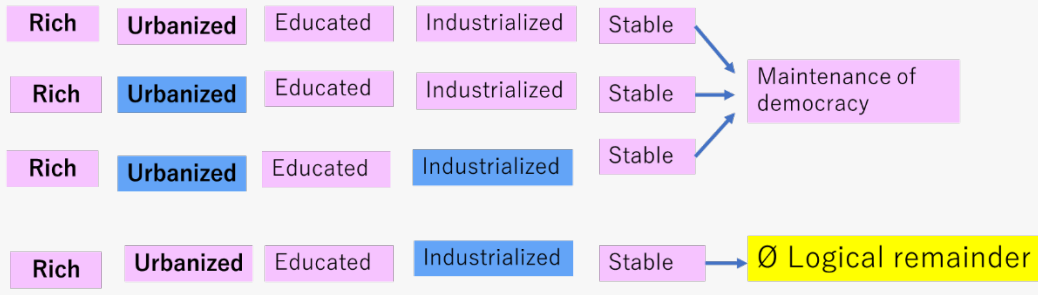
set	condition					result		n	c
	A	B	C	D	E	F	country		
A*B*C*D*E	1	1	1	1	1	1	BEL,CZE,NLD,UK	4	1.00
A*B*C*D*e	1	1	1	1	0	0	GER	1	1.00
A*B*C*d*E	1	1	1	0	1		Logical Remainder		
A*B*C*d*e	1	1	1	0	0		Logical Remainder		
A*B*c*D*E	1	1	0	1	1		Logical Remainder		
A*B*c*D*e	1	1	0	1	0		Logical Remainder		
A*B*c*d*E	1	1	0	0	1		Logical Remainder		
A*B*c*d*e	1	1	0	0	0		Logical Remainder		
A*b*C*D*E	1	0	1	1	1	1	FRA,SWE	2	1.00
A*b*C*D*e	1	0	1	1	0	0	AUT	1	1.00
A*b*C*d*E	1	0	1	0	1	1	FIN,IRL	2	1.00
A*b*C*d*e	1	0	1	0	0		Logical Remainder		
A*b*c*D*E	1	0	0	1	1		Logical Remainder		
A*b*c*D*e	1	0	0	1	0		Logical Remainder		
A*b*c*d*E	1	0	0	0	1		Logical Remainder		
A*b*c*d*e	1	0	0	0	0		Logical Remainder		
a*B*C*D*E	0	1	1	1	1		Logical Remainder		
a*B*C*D*e	0	1	1	1	0		Logical Remainder		
a*B*C*d*E	0	1	1	0	1		Logical Remainder		
a*B*C*d*e	0	1	1	0	0		Logical Remainder		
a*B*c*D*E	0	1	0	1	1		Logical Remainder		
a*B*c*D*e	0	1	0	1	0		Logical Remainder		
a*B*c*d*E	0	1	0	0	1		Logical Remainder		
a*B*c*d*e	0	1	0	0	0		Logical Remainder		
a*b*C*D*E	0	0	1	1	1		Logical Remainder		
a*b*C*D*e	0	0	1	1	0		Logical Remainder		
a*b*C*d*E	0	0	1	0	1	0	EST	1	1.00
a*b*C*d*e	0	0	1	0	0	0	HUN,POL	2	1.00
a*b*c*D*E	0	0	0	1	1		Logical Remainder		
a*b*c*D*e	0	0	0	1	0		Logical Remainder		
a*b*c*d*E	0	0	0	0	1	0	ITA,ROU	2	1.00
a*b*c*d*e	0	0	0	0	0	0	GRC, PRT,ESP	3	1.00

238

minimization

set	A	B	C	D	E	F	country	n	c
$A*B*C*D*E$	1	1	1	1	1	1	BEL,CZE,NLD,UK	4	1.00
$A*b*C*D*E$	1	0	1	1	1	1	FRA,SWE	2	1.00
$A*b*C*d*E$	1	0	1	0	1	1	FIN,IRL	2	1.00

$$(A \wedge B \wedge C \wedge D \wedge E) \vee (A \wedge b \wedge C \wedge D \wedge E) \vee (A \wedge b \wedge C \wedge d \wedge E) \rightarrow R$$



239

240 Fig. 14. Second step of csQCA: minimization of initial solution

241 From Table 10, extracting only the countries that maintain democracy, we can create
 242 the truth table shown in Figure 14. By listing the IDs of sets in the table, we get the
 243 following initial solution:

$$244 \quad A * B * C * D * E + A * b * C * D * E + A * b * C * d * E \subseteq R$$

$$245 \quad (A \wedge B \wedge C \wedge D \wedge E) \vee (A \wedge b \wedge C \wedge D \wedge E) \vee (A \wedge b \wedge C \wedge d \wedge E) \rightarrow R$$

246 Verbalizing this, we get a very lengthy result: “Countries that are wealthy, urbanized,
 247 highly educated, industrialized, and politically stable maintained democracy. Even if
 248 not urbanized, countries that are wealthy, highly educated, industrialized, and
 249 politically stable also maintained democracy. Furthermore, even if not urbanized and
 250 not industrialized, countries that are wealthy, highly educated, and politically stable
 251 maintained democracy.” (Flowchart in Figure 14). Reading this makes us irritated.



252 Immediately upon reading this, one might want to summarize it as
 253 “Wealthy, highly educated, and politically stable countries could maintain democracy”
 254 This is called minimization. The resulting conclusion is referred to as a parsimonious
 255 solution. In essence, it means to explain in a clear and concise manner using the
 256 minimum necessary words. This might be what theorizing is about. Behind this

257 minimization lies an analytical technique called the use of logical remainders for
 258 simplification. We try to simplify the equation obtained directly from the truth table
 259 using Boolean algebra:

$$260 \quad (A \wedge B \wedge C \wedge D \wedge E) \vee (A \wedge b \wedge C \wedge D \wedge E) \vee (A \wedge b \wedge C \wedge d \wedge E) \rightarrow R \quad i$$

261 Using distributive law inversely, we can combine first and second term in left side as
 262 follow.

$$263 \quad (A \wedge C \wedge D \wedge E) \wedge (B + b) \vee (A \wedge b \wedge C \wedge d \wedge E) \rightarrow R \quad ii$$

$$264 \quad (A \wedge C \wedge D \wedge E) \vee (A \wedge b \wedge C \wedge d \wedge E) \rightarrow R \quad iii$$

265 Further minimization is impossible from the logical formula.

266 Fortunately, $A * B * C * d * E$ (the third row of Table 10) is empty set. This means that
 267 $A \wedge B \wedge C \wedge d \wedge E$ is logical remainder.

268 If

$$269 \quad A \wedge B \wedge C \wedge d \wedge E \rightarrow R$$

270 We can add $A \wedge B \wedge C \wedge d \wedge E$ in left side of formula iii and progress the minimization as
 271 follow.

$$272 \quad (A \wedge C \wedge D \wedge E) \vee (A \wedge b \wedge C \wedge d \wedge E) \vee (A \wedge B \wedge C \wedge d \wedge E) \rightarrow R$$

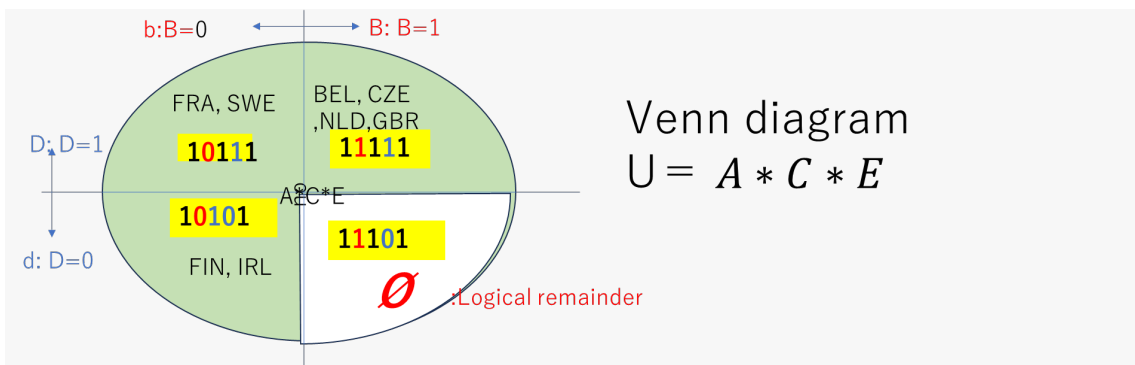
$$273 \quad (A \wedge C \wedge D \wedge E) \vee \{(A \wedge b \wedge C \wedge d \wedge E) \vee (A \wedge B \wedge C \wedge d \wedge E)\} \rightarrow R$$

$$274 \quad (A \wedge C \wedge D \wedge E) \vee \{(A \wedge C \wedge d \wedge E) \wedge (B \vee b)\} \rightarrow R$$

$$275 \quad (A \wedge C \wedge D \wedge E) \vee (A \wedge C \wedge d \wedge E) \rightarrow R$$

$$276 \quad (A \wedge C \wedge E) \wedge (D \vee d) \rightarrow R$$

$$277 \quad A \wedge C \wedge E \rightarrow R$$



278

279

Fig. 15 Venn diagram of $U = A \cap C \cap E$

280 This is the technique for minimization using logical remainder. However, the
 281 explanation should not end here. The moment someone shouted, “Summarize it as
 282 ‘countries that are wealthy, highly educated, and politically stable have maintained
 283 democracy,’” they said, “Ignore things like $A * B * C * d * E$ that don’t exist, and simplify
 284 it by adding it to the logical formula since it’s obviously a country that maintains
 285 democracy.” While this might be true, it is logically reckless.

286 In Figure 15, a Venn diagram is drawn. The entire circular set U in the Venn diagram
 287 is $A * C * E$. This set includes elements with characteristics such as $B * D, B * d, b * D$
 288 and $b * d$. Since there is no intersection (overlap) between them, so they are divided
 289 into four parts and illustrated. The lower right part is an empty set, so it cannot be
 290 colored the same as the others and is left white. The claim of the person angrily
 291 insisting is, “Color it green as a country that maintains democracy.”

292 Whether it should be colored green is a subtle issue. After all, there is no data, and it is
 293 impossible to add data now. One reference is the data of $b * d$ in the third quadrant, to
 294 the left of the white part. The fact that this is green is one reference, and since $b * d$ is
 295 green, there is some basis to think that $B * d$ is also green. It is probably green with a
 296 fairly high probability. In this way, adding logical remainders with evidence,
 297 regardless of logical validity, seems to be a technique of analysis. As evidence, it is also
 298 acceptable to bring in some other examples.

299 There is one problem. In writing this explanation, I referred to Berg-Schlösser D. and
 300 De Meur (1994), which apparently states that logical remainders are easy to use and
 301 convenient, they recommend active use of logical remainder. The reason I say
 302 “apparently” is because I have only read the Japanese translation by Ishida et al.
 303 (2016), 「質的比較分析 (QCA) と関連手法」. I will explain the reason why I did not
 304 read the original separately. Analysts should not cherry-pick convenient logical

305 remainders just because it leads to conclusions that support their claims. This is
306 against research ethics. The assertion that software should be used because it makes
307 this easy is outrageous. It is important to consider the necessity and basis for using
308 those logical remainders. Of course, the burden of proof lies with those who oppose the
309 claim, so one could retort, “If you have complaints, bring counterexamples
310 corresponding to those logical remainders.” However, it has been about 100 years, and
311 there are hardly any similar cases to those in Europe at that time. While falsifiability
312 is a necessary condition for scientific propositions, making claims based on practically
313 unfalsifiable grounds is not commendable and lacks persuasiveness. In fact,
314 simplification can be achieved without overtly using logical remainders. It is a method
315 of evaluating logical consistency.

316 The method is shown in Figure 16. Among the three tables in Figure 16, the top table
317 examines $A \wedge C \rightarrow R$. First, $A \wedge C$ meaning selecting cases where both A and C are 1 (in
318 the Excel sheet, prioritize A and C and sort them in descending order). Count the
319 number of countries that meet this condition, dividing them into countries that
320 maintain democracy and those that collapse. Using the following formula, calculate the
321 proportion of countries that maintain democracy relative to the total number of
322 countries:

$$323 \quad c = \frac{\text{total number of countries maintained democracy}}{\text{total numbers of countries belonging the conditon}}$$

324 This proportion represents the consistency of the claim when $A \wedge C \rightarrow R$. Since csQCA
325 is crispy and does not accept intermediate values other than 1, combinations with a
326 consistency other than 1 are not accepted (rejected). The consistency of $A \wedge C \rightarrow R$ is
327 0.80, so it is rejected. Similarly, $C \wedge E \rightarrow R$ is also rejected with a consistency of 0.89.
328 The only accepted combination is $A \wedge E \rightarrow R$ with a consistency of 1. Although it might
329 be obvious without calculation, the method of quantifying and comparing consistency is
330 central to fsQCA (csQCA can be considered a special case of fsQCA).

sets ID	condition					Result		
	A	B	C	D	E	R	country	n
A*B*C*D*E	1	1	1	1	1	1	BEL,CZE,NLD,UK	4
A*B*C*d*E	1	1	1	0	1	0	Logical Remainder	
A*b*C*D*E	1	0	1	1	1	1	FRA,SWE	2
A*b*C*d*E	1	0	1	0	1	1	FIN,IRL	2
A*B*C*D*e	1	1	1	1	0	0	0 GER	1
A*B*C*d*e	1	1	1	0	0	0	Logical Remainder	
A*b*C*D*e	1	0	1	1	0	0	0 AUT	1
A*b*C*d*e	1	0	1	0	0	0	Logical Remainder	

331

$$\begin{aligned}
 & A * C * E + A * C * e \\
 &= A * C * (E + e) \\
 &= A * C \\
 & \text{consistency}(A \wedge C \rightarrow R) \\
 & c = \frac{4 + 2 + 2}{4 + 2 + 2 + 1 + 1} = 0.80
 \end{aligned}$$

Reject

sets ID	condition					Result		
	A	B	C	D	E	R	country	n
A*B*C*D*E	1	1	1	1	1	1	BEL,CZE,NLD,UK	4
A*B*C*d*E	1	1	1	0	1	0	Logical Remainder	
A*b*C*D*E	1	0	1	1	1	1	FRA,SWE	2
A*b*C*d*E	1	0	1	0	1	1	FIN,IRL	2
a*B*C*D*E	0	1	1	1	1	0	Logical Remainder	
a*B*C*d*E	0	1	1	0	1	0	Logical Remainder	
a*b*C*D*E	0	0	1	1	1	0	Logical Remainder	
a*b*C*d*E	0	0	1	0	1	0	0 EST	1

332

$$\begin{aligned}
 & A * C * E + a * C * E \\
 &= C * E * (A + a) \\
 &= C * E \\
 & \text{consistency}(C \wedge E \rightarrow R) \\
 & c = \frac{4 + 2 + 2}{4 + 2 + 2 + 1} = 0.89
 \end{aligned}$$

Reject

sets ID	condition					Result		
	A	B	C	D	E	R	country	n
A*B*C*D*E	1	1	1	1	1	1	BEL,CZE,NLD,UK	4
A*B*C*d*E	1	1	1	0	1	0	Logical Remainder	
A*B*c*D*E	1	1	0	1	1	0	Logical Remainder	
A*B*C*d*E	1	1	0	0	1	0	Logical Remainder	
A*b*C*D*E	1	0	1	1	1	1	FRA,SWE	2
A*b*C*d*E	1	0	1	0	1	1	FIN, FIN	2
A*b*c*D*E	1	0	0	1	1	0	Logical Remainder	
A*b*c*d*E	1	0	0	0	1	0	Logical Remainder	

333

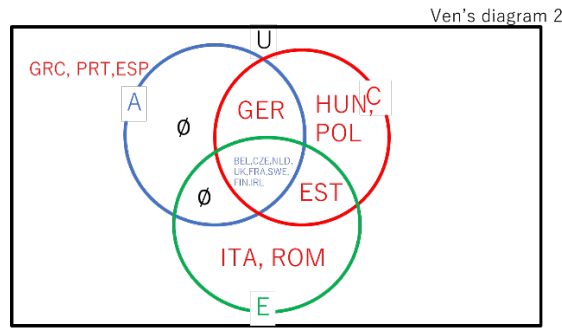
$$\begin{aligned}
 & A * C * E + A * c * E \\
 &= A * E * (C + c) \\
 &= A * E \\
 & \text{consistency}(A \wedge E \rightarrow R) \\
 & c = \frac{4 + 2 + 2}{4 + 2 + 2} = 1.00
 \end{aligned}$$

Accept

334

Fig 16. Minimizing of condition by consistency

335



336

337

Fig. 17 Venn diagram of A, C, E

$A * C * D * e$

set	condition					result	n	c	
	A	B	C	D	E	R			country
$A^*B^*C^*D^*e$	1	1	1	1	0	0	GER	1	1.00
$A^*b^*C^*D^*e$	1	0	1	1	0	0	AUT	1	1.00

$a * b * d$

set	condition					result	n	c	
	A	B	C	D	E	res			country
$a^*b^*C^*d^*E$	0	0	1	0	1	0	EST	1	1.00
$a^*b^*C^*d^*e$	0	0	1	0	0	0	HUN,POL	2	1.00
$a^*b^*c^*d^*E$	0	0	0	0	1	0	ITA,ROU	2	1.00
$a^*b^*c^*d^*e$	0	0	0	0	0	0	GRC, PRT,ESP	3	1.00

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Fig. 18. Conditions for collapse of democracy, from initial solution to intermediate solution

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Figure 17 shows the relationship of sets A, C, and E in a Venn diagram. Since $A \cap C \cap E$ is a subset of $A \cap E$, when $A \wedge E \rightarrow R$ holds, $A \wedge C \wedge E \rightarrow R$ also holds. The explanation based on consistency is more convincing than an explanation with logical remainders. $A \cap E$ includes $A \cap c \cap E$ as a subset. Since $A \cap c \cap E$ is an empty set (\emptyset), $A \wedge c \wedge E$ is a logical remainder. In fact, even in simplification by confirming consistency, there is an implicit interpretation of logical remainders. In other words, it can be overturned if there is a counterexample. It is confirmed that $A \wedge C \wedge E \rightarrow R$ has no counterexamples in the data.

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Next, we will examine the conditions that lead to the collapse of democracy. The condition is represented as $R=0$ or r . Returning to the truth table of conditions in Table 10, we create a list of countries where democracy has collapsed, as shown in the left table of Figure 18. This is the initial solution. These can be divided into two groups.

353 One group consists of countries with A=1, C=1, D=1, and E=0, represented by the
 354 condition $A \wedge C \wedge D \wedge e$, which includes Germany and Austria. The other group consists
 355 of countries with A=0, B=0, and D=0, represented by the condition $a \wedge b \wedge d$ which
 356 includes Estonia, Hungary, Poland, Italy, Romania, Greece, Portugal, and Spain. Thus,
 357 the intermediate solution is:

358
$$(A \wedge C \wedge D \wedge e) \vee (a \wedge b \wedge d) \rightarrow r$$

359 We then simplify each of the two terms in this intermediate solution to find the most
 360 parsimonious solution. The process of simplifying $A \wedge C \wedge D \wedge e$ from aggregation to

$A \rightarrow r$							$r: R=0$							$C \rightarrow r$							$r: R=0$														
set	A	B	C	D	E	R	country	n	set	A	B	C	D	E	R	country	n	set	A	B	C	D	E	R	country	n	set	A	B	C	D	E	R	country	n
A*B*C*D*e	1	1	1	1	0	0	GER	1	A*B*C*D*e	1	1	1	1	0	0	GER	1	A*B*C*D*e	1	1	1	1	0	0	GER	1	A*B*C*D*e	1	1	1	1	0	0	GER	1
A*B*C*d*e	1	1	1	0	0		Logical Remainder		A*B*C*d*e	1	1	1	0	0		Logical Remainder		A*B*C*d*e	1	1	1	0	0		Logical Remainder		A*B*C*d*e	1	1	1	0	0		Logical Remainder	
A*B*c*D*e	1	1	0	1	0		Logical Remainder		A*B*c*D*e	1	1	0	1	0		AUT	1	A*B*c*D*e	1	1	0	1	0		AUT	1	A*B*c*D*e	1	1	0	1	0		Logical Remainder	
A*B*c*d*e	1	1	0	0	0		Logical Remainder		A*B*c*d*e	1	1	0	0	0		Logical Remainder		A*B*c*d*e	1	1	0	0	0		Logical Remainder		A*B*c*d*e	1	1	0	0	0		Logical Remainder	
A*b*C*D*e	1	0	1	1	0	0	AUT	1	A*b*C*D*e	1	0	1	1	1	1	BEL,CZE,NLD,UK	4	A*b*C*D*e	1	0	1	1	1	1	BEL,CZE,NLD,UK	4	A*b*C*D*e	1	0	1	1	1	1	BEL,CZE,NLD,UK	4
A*b*C*d*e	1	0	1	0	0		Logical Remainder		A*b*C*d*e	1	0	1	0	1		Logical Remainder		A*b*C*d*e	1	0	1	0	1		Logical Remainder		A*b*C*d*e	1	0	1	0	1		Logical Remainder	
A*b*c*D*e	1	0	0	1	0		Logical Remainder		A*b*c*D*e	1	0	0	1	1	1	FRA,SWE	2	A*b*c*D*e	1	0	0	1	1	1	FRA,SWE	2	A*b*c*D*e	1	0	0	1	1	1	FRA,SWE	2
A*b*c*d*e	1	0	0	0	0		Logical Remainder		A*b*c*d*e	1	0	0	0	1	1	FIN,IRL	2	A*b*c*d*e	1	0	0	0	1	1	FIN,IRL	2	A*b*c*d*e	1	0	0	0	1	1	FIN,IRL	2
A*B*C*D*E	1	1	1	1	1	1	BEL,CZE,NLD,UK	4	A*B*C*D*E	1	1	1	1	1	1	BEL,CZE,NLD,UK	4	A*B*C*D*E	1	1	1	1	1	1	BEL,CZE,NLD,UK	4	A*B*C*D*E	1	1	1	1	1	1	BEL,CZE,NLD,UK	4
A*B*C*d*E	1	1	1	0	1		Logical Remainder		A*B*C*d*E	1	1	1	0	1		Logical Remainder		A*B*C*d*E	1	1	1	0	1		Logical Remainder		A*B*C*d*E	1	1	1	0	1		Logical Remainder	
A*B*c*D*E	1	1	0	1	1		Logical Remainder		A*B*c*D*E	1	1	0	1	1		Logical Remainder		A*B*c*D*E	1	1	0	1	1		Logical Remainder		A*B*c*D*E	1	1	0	1	1		Logical Remainder	
A*B*c*d*E	1	1	0	0	1		Logical Remainder		A*B*c*d*E	1	1	0	0	1		Logical Remainder		A*B*c*d*E	1	1	0	0	1		Logical Remainder		A*B*c*d*E	1	1	0	0	1		Logical Remainder	
A*b*C*D*E	1	0	1	1	1	1	FRA,SWE	2	A*b*C*D*E	1	0	1	1	1	1	FRA,SWE	2	A*b*C*D*E	1	0	1	1	1	1	FRA,SWE	2	A*b*C*D*E	1	0	1	1	1	1	FRA,SWE	2
A*b*C*d*E	1	0	1	0	1	1	FIN,IRL	2	A*b*C*d*E	1	0	1	0	1	1	FIN,IRL	2	A*b*C*d*E	1	0	1	0	1	1	FIN,IRL	2	A*b*C*d*E	1	0	1	0	1	1	FIN,IRL	2
A*b*c*D*E	1	0	0	1	1		Logical Remainder		A*b*c*D*E	1	0	0	1	1		Logical Remainder		A*b*c*D*E	1	0	0	1	1		Logical Remainder		A*b*c*D*E	1	0	0	1	1		Logical Remainder	
A*b*c*d*E	1	0	0	0	1		Logical Remainder		A*b*c*d*E	1	0	0	0	1		Logical Remainder		A*b*c*d*E	1	0	0	0	1		Logical Remainder		A*b*c*d*E	1	0	0	0	1		Logical Remainder	

361
$$c = \frac{1 + 1}{1 + 1 + 4 + 2 + 2} = \frac{2}{10} = 0.20$$

$$c = \frac{1 + 1 + 2 + 1}{1 + 1 + 4 + 2 + 2 + 2 + 1} = \frac{5}{13} = 0.38$$

$D \rightarrow r$							$r: R=0$							$e \rightarrow r$							$r: R=0$														
set	A	B	C	D	E	R	country	n	set	A	B	C	D	E	R	country	n	set	A	B	C	D	E	R	country	n	set	A	B	C	D	E	R	country	n
A*B*C*D*e	1	1	1	1	0	0	GER	1	A*B*C*D*e	1	1	1	1	0	0	GER	1	A*B*C*D*e	1	1	1	1	0	0	GER	1	A*B*C*D*e	1	1	1	1	0	0	GER	1
A*b*C*D*e	1	0	1	1	0	0	AUT	1	A*b*C*D*e	1	0	1	1	0	0	AUT	1	A*b*C*D*e	1	0	1	1	0	0	AUT	1	A*b*C*D*e	1	0	1	1	0	0	AUT	1
A*B*C*D*E	1	1	1	1	1	1	BEL,CZE,NLD,UK	4	A*B*C*D*E	1	0	1	1	1	1	FRA,SWE	2	A*B*C*D*E	1	0	1	1	1	1	FRA,SWE	2	A*B*C*D*E	1	0	1	1	1	1	FRA,SWE	2
A*b*C*D*E	1	0	1	1	1		Logical Remainder		A*b*C*D*E	1	0	1	1	0		Logical Remainder		A*b*C*D*E	1	0	1	1	0		Logical Remainder		A*b*C*D*E	1	0	1	1	0		Logical Remainder	
A*b*C*d*E	1	0	1	0	0		Logical Remainder		A*b*C*d*E	1	0	1	0	0		Logical Remainder		A*b*C*d*E	1	0	1	0	0		Logical Remainder		A*b*C*d*E	1	0	1	0	0		Logical Remainder	
A*b*c*D*E	1	0	0	1	1		Logical Remainder		A*b*c*D*E	1	0	0	1	1		Logical Remainder		A*b*c*D*E	1	0	0	1	1		Logical Remainder		A*b*c*D*E	1	0	0	1	1		Logical Remainder	
A*B*C*D*E	0	1	1	1	1		Logical Remainder		A*B*C*D*E	0	1	1	1	0		Logical Remainder		A*B*C*D*E	0	1	1	1	0		Logical Remainder		A*B*C*D*E	0	1	1	1	0		Logical Remainder	
A*b*C*D*E	0	0	1	1	1		Logical Remainder		A*b*C*D*E	0	0	1	1	0		Logical Remainder		A*b*C*D*E	0	0	1	1	0		Logical Remainder		A*b*C*D*E	0	0	1	1	0		Logical Remainder	
A*B*c*D*E	1	1	0	1	0		Logical Remainder		A*B*c*D*E	1	1	0	1	0		Logical Remainder		A*B*c*D*E	1	1	0	1	0		Logical Remainder		A*B*c*D*E	1	1	0	1	0		Logical Remainder	
A*b*c*D*E	1	0	0	1	0		Logical Remainder		A*b*c*D*E	1	0	0	1	0		Logical Remainder		A*b*c*D*E	1	0	0	1	0		Logical Remainder		A*b*c*D*E	1	0	0	1	0		Logical Remainder	
A*B*c*D*E	1	1	0	1	1		Logical Remainder		A*B*c*D*E	1	1	0	1	1		Logical Remainder		A*B*c*D*E	1	1	0	1	1		Logical Remainder		A*B*c*D*E	1	1	0	1	1		Logical Remainder	
A*b*c*D*E	1	0	0	1	1		Logical Remainder		A*b*c*D*E	1	0	0	1	1		Logical Remainder		A*b*c*D*E	1	0	0	1	1		Logical Remainder		A*b*c*D*E	1	0	0	1	1		Logical Remainder	
A*b*c*d*E	1	0	0	0	1		Logical Remainder		A*b*c*d*E	1	0	0	0	1		Logical Remainder		A*b*c*d*E	1	0	0	0	1		Logical Remainder		A*b*c*d*E	1	0	0	0	1		Logical Remainder	
A*B*c*D*E	0	1	0	1	1		Logical Remainder		A*B*c*D*E	0	1	0	1	0		Logical Remainder		A*B*c*D*E	0	1	0	1	0		Logical Remainder		A*B*c*D*E	0	1	0	1	0		Logical Remainder	
A*b*c*D*E	0	0	0	1	0		Logical Remainder		A*b*c*D*E	0	0	0	1	0		Logical Remainder		A*b*c*D*E	0	0	0	1	0		Logical Remainder		A*b*c*D*E	0	0	0	1	0		Logical Remainder	
A*B*c*D*E	0	1	0	1	1		Logical Remainder		A*B*c*D*E	0	1	0	1	0		Logical Remainder		A*B*c*D*E	0	1	0	1	0		Logical Remainder		A*B*c*D*E	0	1	0	1	0		Logical Remainder	
A*b*c*D*E	0	0	0	1	0		Logical Remainder		A*b*c*D*E	0	0	0	1	0		Logical Remainder		A*b*c*D*E	0	0	0	1	0		Logical Remainder		A*b*c*D*E	0	0	0	1	0		Logical Remainder	
A*B*c*D*E	0	1	0	1	1		Logical Remainder		A*B*c*D*E	0	1	0	1	0		Logical Remainder		A*B*c*D*E	0	1	0	1	0		Logical Remainder		A*B*c*D*E	0	1	0	1	0		Logical Remainder	
A*b*c*D*E	0	0	0	1	1		Logical Remainder		A*b*c*D*E	0	0	0	1	1		Logical Remainder		A*b*c*D*E	0	0	0	1	1		Logical Remainder		A*b*c*D*E	0	0	0	1	1		Logical Remainder	
A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder	
A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder	
A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder	
A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder	
A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder	
A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder	
A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder	
A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder	
A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder	
A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder	
A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder	
A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0	1		Logical Remainder		A*b*c*d*E	0	0	0	0				

364 consistency calculation, is shown in Figure 19. We examined $A \rightarrow r$, $C \rightarrow r$, $D \rightarrow r$, and
 365 $e \rightarrow r$. Predictably, $A \rightarrow r$, $C \rightarrow r$, and $D \rightarrow r$ are not possible. In fact, their consistency is
 366 extremely low. On the other hand, $e \rightarrow r$ has a consistency of 1 and can be adopted as
 367 the most parsimonious solution.

368 Next, the process from intermediate solution to the most parsimonious solution for $a \wedge$
 369 $b \wedge d \rightarrow r$ is shown in Figure 20. This development yields the most parsimonious
 370 solution:

set	A	B	C	D	E	R	country	
a*B*C*D*e	0	1	1	1	0		Logical Remainder	
a*b*C*D*e	0	0	1	1	0		Logical Remainder	
a*B*c*D*e	0	1	0	1	0		Logical Remainder	
a*b*c*D*e	0	0	0	1	0		Logical Remainder	
a*B*C*d*e	0	1	1	0	0		Logical Remainder	
a*b*c*d*e	0	0	1	0	0	0	HUN,POL	2
a*B*c*d*e	0	1	0	0	0		Logical Remainder	
a*b*c*d*e	0	0	0	0	0	0	GRC, PRT,ESP	2
a*B*C*D*E	0	1	1	1	1		Logical Remainder	
a*b*c*D*E	0	0	1	1	1		Logical Remainder	
a*B*c*D*E	0	1	0	1	1		Logical Remainder	
a*b*c*D*E	0	0	0	1	1		Logical Remainder	
a*B*C*d*E	0	1	1	0	1		Logical Remainder	
a*b*c*d*E	0	0	1	0	1	0	EST	1
a*B*c*d*E	0	1	0	0	1		Logical Remainder	
a*b*c*d*E	0	0	0	0	1	0	ITA,ROU	2

371 $c = \frac{2 + 2 + 1 + 2}{2 + 2 + 1 + 2} = \frac{7}{7} = 1.00$ Accept

set	A	B	C	D	E	R	country	n
a*b*C*D*e	0	0	1	1	0		Logical Remainder	
a*b*c*D*e	0	0	0	1	0		Logical Remainder	
a*b*c*d*e	0	0	1	0	0	0	HUN,POL	2
a*b*c*d*e	0	0	0	0	0	0	GRC, PRT,ESP	3
a*b*c*D*E	0	0	1	1	1		Logical Remainder	
a*b*c*d*E	0	0	0	1	1		Logical Remainder	
a*b*c*d*E	0	0	1	0	1	0	EST	1
a*b*c*d*E	0	0	0	0	1	0	ITA,ROU	2
A*b*C*D*E	1	0	1	1	0	0	AUT	1
A*b*c*D*E	1	0	0	1	0		Logical Remainder	
A*b*c*d*E	1	0	1	0	0		Logical Remainder	
A*b*c*d*E	1	0	0	0	0		Logical Remainder	
A*b*C*D*E	1	0	1	1	1	1	FRA,SWE	2
A*b*c*D*E	1	0	0	1	1		Logical Remainder	
A*b*c*d*E	1	0	1	0	1	1	FIN,IRL	2
A*b*c*d*E	1	0	0	0	1		Logical Remainder	

$c = \frac{2 + 3 + 1 + 2 + 1}{2 + 3 + 1 + 2 + 1 + 2 + 2} = \frac{9}{13} = 0.69$

$d \rightarrow r$

set	A	B	C	D	E	R	country	
a*b*C*d*e	0	0	1	0	0	0	HUN,POL	2
a*b*c*d*e	0	0	0	0	0	0	GRC, PRT,ESP	3
a*b*c*d*E	0	0	1	0	1	0	EST	1
a*b*c*d*E	0	0	0	0	1	0	ITA,ROU	2
A*b*C*d*e	1	0	1	0	0		Logical Remainder	
A*b*c*d*e	1	0	0	0	0		Logical Remainder	
A*b*C*d*E	1	0	1	0	1	1	FIN,IRL	2
A*b*c*d*E	1	0	0	0	1		Logical Remainder	
a*B*C*d*e	0	1	1	0	0		Logical Remainder	
a*B*c*d*e	0	1	0	0	0		Logical Remainder	
a*B*C*d*E	0	1	1	0	1		Logical Remainder	
a*B*c*d*E	0	1	0	0	1		Logical Remainder	
A*B*C*d*e	1	1	1	0	0		Logical Remainder	
A*B*c*d*e	1	1	0	0	0		Logical Remainder	
A*B*C*d*E	1	1	1	0	1		Logical Remainder	
A*B*c*d*E	1	1	0	0	1		Logical Remainder	

372 $c = \frac{2 + 3 + 1 + 2}{2 + 3 + 1 + 2 + 2} = \frac{8}{10} = 0.80$

373 Fig. 20. Intermediate solution $a \wedge b \wedge d \rightarrow r$ to final solution

374 $a \rightarrow r$

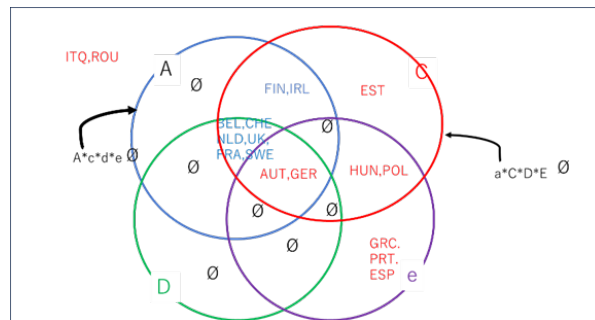
375 The logical OR of the two parsimonious solutions is:

376 $a \vee e \rightarrow r$

377 In everyday language, this means that if there is “poverty” or “political instability,”
378 democracy will collapse. If either “poverty” or “political instability” exists, democracy
379 will collapse.

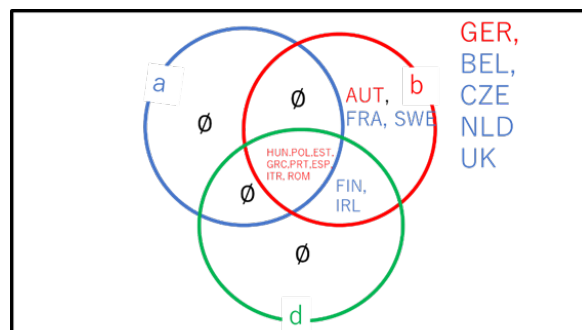
380 Verify the plausibility of this proposition using a Venn diagram. Figure 21 is a Venn
381 diagram of A, C, D, and e, while Figure 22 is a Venn diagram of a, b, and d.

382 Theoretically, 4 simple conditions make of $2^4 = 16$ combinations of conditions subsets,
383 but in Figure 21, $A * c * d * e$ and $a * C * D * E$ are hidden behind other sets. Both $A * c * d * e$
384 and $a * C * D * E$ are empty sets, and $A * c * d * e$ belongs to e ($A * c * d * e \subset e$). e
385 consists of eight subsets, and seven countries (HUN, POL, AUT, GER, GRC, PRT, ESP)



386

387 Fig. 21 Venn diagram of A, B,C, and e



388

389 Fig. 21 Venn diagram of a, b, and d

390 belong to e . All of these are countries where democracy has collapsed. There is no
 391 logical problem in concluding $e \rightarrow r$ meaning “if politics is unstable, democracy will
 392 collapse,” as the most parsimonious solution, but it should be noted that five out of the
 393 eight subsets are empty sets. Additionally, e includes seven countries where democracy
 394 has collapsed. There are ten countries where democracy has collapsed. When s
 395 concluded, it explains 70% of the countries. In contrast, concluding $e \wedge C \rightarrow r$ explains
 396 40% of the total, and $e \wedge D \rightarrow r$ explains only 20%. Even $e \wedge C \wedge D \rightarrow R$ explains only
 397 20%. If that is true, which part of the result to emphasize in the conclusion should be
 398 left to the analyst’s judgment, but in general cases without special analytical purposes,
 399 it is common sense to conclude what applies to a wider range. The ratio of how much of
 400 the result can be explained is called coverage. Comparing the coverage, the final
 401 solution is concluded as $e \rightarrow r$. Looking at Figure 22, a includes eight countries where
 402 democracy has collapsed, and no countries where democracy is maintained. Most
 403 parsimonious solution is $a \rightarrow r$.

404 From the perspective of consistency, this is also logically unproblematic. However,
 405 including logical residues, $a \wedge b \rightarrow r$ and $a \wedge d \rightarrow r$ also hold. Originally, $a \wedge b \wedge d \rightarrow r$
 406 was established, so the final conclusion is chosen from the four solutions. In this
 407 case, $a \rightarrow r$, $a \wedge b \rightarrow r$, $a \wedge d \rightarrow r$, and $a \wedge b \wedge d \rightarrow r$ are established, and the coverage
 408 is the same. Explaining phenomena with as few factors as possible might be one of the
 409 principles of science. If so, the general idea would be to conclude $a \rightarrow r$, with a as the
 410 core condition a as the core condition, and b and d as peripheral conditions, with $a \rightarrow r$
 411 as the final conclusion.

412 In conclusion, it becomes:

413
$$a \vee e \rightarrow r$$

414 Let’s write this conclusion alongside the conclusion for countries maintaining
 415 democracy:

416
$$A \wedge E \rightarrow R, \quad a \vee e \rightarrow r$$

417 Have you noticed? It follows De Morgan’s laws. De Morgan’s laws state:

418
$$\widetilde{A * B} = \widetilde{A} + \widetilde{B}$$

419 It might be a bit confusing, so let’s write it properly using set notation:

420 $A \cap E \subseteq R, \quad \tilde{A} \cup \tilde{E} \subseteq \tilde{R}$

421 De Morgan's law states:

422 $\tilde{A} \cup \tilde{B} = \widetilde{A \cap B}$

423 In other words, the negation of $(A \wedge E)$ will always be the negation of (R) . This means:

424 $A \wedge E \Leftrightarrow R$

425 $A \wedge E$ and R are equivalent, being necessary and sufficient conditions for each other.

426 In everyday language, if a country is prosperous and politically stable, democracy will
427 always be maintained. Conversely, if these conditions are not met, democracy will
428 inevitably collapse.

429 What I am trying to say here is not that this will always be the case in csQCA. Nor am
430 I suggesting that you should aim for this in your analysis. Such situations are rare and
431 feel quite unnatural. I suspect that the original data might have been constructed to
432 produce these results. After all, there is no explanation of how the degree of democracy
433 maintenance was evaluated. Once you start doubting, there's no end to it. However, I
434 won't delve deeper into this. The purpose of this explanation is to ultimately
435 demonstrate what can be done with fsQCA2 and how to do it. In the next chapter, I will
436 explain fsQCA.

437 The reference book I am using explains mvQCA, which involves dividing the data into
438 three or more categories instead of binary values for csQCA. This means that for some
439 items in the original data, instead of categorizing them as simply large or small, they
440 are divided into large, medium, and small, or even more detailed categories. After that,
441 you just need to write the data divided into three categories in the truth table as 0, 1,
442 and 2. Nothing else changes. If necessary, I can provide an explanation, but for now, I
443 will explain fsQCA.

444